# DIMENSION ESTIMATE FOR A SET OBTAINED FROM A THREE-DIMENSIONAL NON-PERIODIC SELF-AFFINE TILING 

Taizo Sadahiro

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#### Abstract

Some computational results on the three-dimensional Pisot tiling generated by the roots of $x^{4}-x^{3}-x^{2}-x-1=0$ are shown. We show an upper bound of the Hausdorff dimension of the set which is a projection of the intersection of three-tiles to the plane.


1 Introduction A tile is a compact subset of $\mathbf{R}^{n}$ which is equal to the closure of its interior. A set of tiles $\mathcal{T}$ is a tiling of $\mathbf{R}^{n}$, if $\mathcal{T}$ is a covering of $\mathbf{R}^{n}$ such that the intersection of interiors of any two tiles in $\mathcal{T}$ is empty. A tiling $\mathcal{T}$ is called self-affine, if there is an affine map such that the image of a tile in $\mathcal{T}$ is a union of tiles of $\mathcal{T}$.

Self-affine tilings are of special interest because of their relation to several topics of recent research, Markov partitions for toral automorphisms [11, 4, 1], wavelet theory [7, 10] and real quasi-crystal [8, 13]. There are numerous studies on self-affine tilings. All of the examples used there are about one or two-dimensional cases. Explicit examples of higher than or equal to three-dimensional cases have not been studied.

In this paper, we show some explicit computational results on the 3 -dimensional nonperiodic self-affine tiling generated by the roots of the equation $x^{4}-x^{3}-x^{2}-x-1=0$. We construct a Mauldin-Williams graph [9] of intersection of three tiles. The tiling we treat here is a Pisot tiling. Here we do not give the precise and general definition of the Pisot tilings. For further details, see $[2,15,14,3]$. The tiling is constructed as follows: The polynomial $x^{4}-x^{3}-x^{2}-x-1$ is irreducible over $\mathbf{Q}$ and has four distinct roots,

$$
\begin{aligned}
\beta & =1.927561975482925304261905 \cdots \\
\gamma & =-0.7748041132154338540924032 \cdots \\
\alpha & =-0.07637893113374572508475 \cdots-0.8147036471703865268416 \cdots i \\
\bar{\alpha} & =-0.07637893113374572508475 \cdots+0.8147036471703865268416 \cdots i
\end{aligned}
$$

So $\beta$ is a Pisot number, that is, $\beta$ is an algebraic integer greater than 1 and all of its Galois conjugates over $\mathbf{Q}$ are strictly inside the unit circle. Let $w=d_{-l} d_{-l+1} \cdots d_{-1}$ be a word over $\{0,1\}$. A tile $T(w)$ is defined as follows:

$$
T(w)=\left\{\left(\sum_{i=-l}^{\infty} a_{i} \alpha^{i}, \sum_{i=-l}^{\infty} a_{i} \gamma^{i}\right): \begin{array}{l}
a_{i} \in\{0,1\}, a_{i} \times a_{i+1} \times a_{i+2} \times a_{i+3}=0 \\
a_{-l} a_{-l+1} \cdots a_{-1}=w
\end{array}\right\}
$$

which is a subset of $\mathbf{C} \times \mathbf{R} \simeq \mathbf{R}^{3}$. The tiling $\mathcal{T}$ is defined by

$$
\mathcal{T}:=\left\{T(w): w \in\{0,1\}^{*}\right\}
$$

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where $\{0,1\}^{*}$ denotes the set of all of the words over $\{0,1\}$ (including the empty word $\epsilon$ ).
An automaton is a (directed labeled) graph. For an automaton $M$, we denote by $\mathcal{V}(M)$ the vertex set of $M$ and by $\mathcal{E}(M)$ the edge set of $M$. Every edge $e \in \mathcal{E}(M)$ has the starting point $s(e) \in \mathcal{V}(M)$ and the end point $t(e) \in \mathcal{V}(M)$, and carry a label $l(e) \in \Sigma(M)$ where $\Sigma(M)$ is a finite set called the alphabet of $M$. A sequence of edges $e_{1} \cdots e_{l}$ is called a path of $M$ if $t\left(e_{i}\right)=s\left(e_{i+1}\right)$. An automaton has a special vertex $i_{M}$ called the initial state of $M . \Sigma^{*}$ denotes the set of all of the words over an alphabet $\Sigma$. A word $w=a_{1} \cdots a_{l} \in \Sigma^{*}$ is accepted by $M$ if there exists a path $p=e_{1} \cdots e_{l}$ starting from $i_{M}$ such that $l\left(e_{1}\right) \cdots l\left(e_{l}\right)=w$. An infinite word $\left(a_{i}\right)_{i \geq 0}$ over $\Sigma$ is accepted by $M$ if $a_{0} \cdots a_{h}$ is accepted by $M$ for all $h \geq 0$. We denote by $L(M, i)$ the set of infinite words accepted by $M$ with the initial state $i$.

Example 1 The automaton shown in Figure 1 accepts the words over $\{0,1\}$ which does not include 11 as a subword.


Figure 1:

The intersection of tiles are determined by automata. See [12] for the proof of the following theorem.

Theorem 1 (Sadahiro) The intersection of tiles are represented by automata: For any $n$ tiles $T\left(w_{1}\right), T\left(w_{2}\right), \ldots, T\left(w_{n}\right)$, there exists an automaton $m$ with the following property.

$$
\left(\sum_{i=-l}^{\infty} a_{i} \alpha^{i}, \sum_{i=-l}^{\infty} a_{i} \gamma^{i}\right) \in T\left(w_{1}\right) \cap T\left(w_{2}\right) \cap \cdots \cap T\left(w_{n}\right)
$$

if and only if $a_{-l} \cdots a_{-1}=w_{1}$ and $a_{-l} a_{-l+1} \cdots a_{h}$ is accepted by $m$ for any $h \geq 0$.
An infinite word accepted by the automaton in the theorem above determines a point in the intersection. For example, the automaton which represents $T(0) \cap T(1) \cap T(11) \cap T(111)$ is shown in Figure 2, from which we can see $T(0) \cap T(1) \cap T(11) \cap T(111)$ consists of only one point $(-1,-1) \in \mathbf{C} \times \mathbf{R}$.

In fact, the following four presentations of $(-1,-1)$ exist:

$$
\begin{aligned}
(-1,-1) & =\left(\sum_{n=0}^{\infty} \alpha^{4 n}\left(\alpha+\alpha^{2}+\alpha^{3}\right), \sum_{n=0}^{\infty} \gamma^{4 n}\left(\gamma+\gamma^{2}+\gamma^{3}\right)\right)(\in T(0)) \\
& =\left(\frac{1}{\alpha}+\sum_{n=0}^{\infty} \alpha^{4 n}\left(\alpha+\alpha^{2}+\alpha^{4}\right), \frac{1}{\gamma}+\sum_{n=0}^{\infty} \gamma^{4 n}\left(\gamma+\gamma^{2}+\gamma^{4}\right)\right)(\in T(1))
\end{aligned}
$$



Figure 2:

$$
\begin{aligned}
& =\left(\frac{1}{\alpha^{2}}+\frac{1}{\alpha}+\sum_{n=0}^{\infty} \alpha^{4 n}\left(\alpha+\alpha^{3}+\alpha^{4}\right),\right. \\
& \\
& =\left(\frac{1}{\gamma^{2}}+\frac{1}{\gamma}+\sum_{n=0}^{\infty} \gamma^{4 n}\left(\gamma+\gamma^{3}+\gamma^{4}\right)\right)(\in T(11)) \\
& \alpha^{3}+\frac{1}{\alpha^{2}}+\frac{1}{\alpha}+\sum_{n=0}^{\infty} \alpha^{4 n}\left(\alpha^{2}+\alpha^{3}+\alpha^{4}\right), \\
& \\
& \left.\frac{1}{\gamma^{3}}+\frac{1}{\gamma^{2}}+\frac{1}{\gamma}+\sum_{n=0}^{\infty} \gamma^{4 n}\left(\gamma^{2}+\gamma^{3}+\gamma^{4}\right)\right)(\in T(111)) .
\end{aligned}
$$

2 Dimension of $T(0) \cap T(1) \cap T(11)$ We will study the following set in $\mathbf{C}$ :

$$
E=\left\{\sum_{i=0}^{\infty} a_{i} \alpha^{i}:\left(\sum_{i=0}^{\infty} a_{i} \alpha^{i}, \sum_{i=0}^{\infty} a_{i} \gamma^{i}\right) \in T(0) \cap T(1) \cap T(11)\right\} .
$$

Figure 3 shows $E$. The automaton which accepts words determining points in $T(0) \cap T(1) \cap$ $T(11)$ is shown in the appendix.

A cycle is a directed graph $H$ for which there is a closed path which passes into every vertex exactly once and such that every edge of $H$ is an edge of this path. A directed graph $H$ is strongly connected provided that whenever each of $x$ and $y$ is a vertex of $H$, then there is a path from $x$ to $y$.

A strongly connected component of $G$ is a maximal subgraph $H$ of $G$ such that $H$ is strongly connected. It is clear that the strongly connected components of $G$ are pairwise disjoint. A vertex is not considered to be strongly connected unless it is looped on itself.

The automaton which represents $T(0) \cap T(1) \cap T(11)$ is decomposed into strongly connected components as is shown in Figure 4. Every strongly connected components except a special component $X$ consists of one cycle. Figure 5 shows the component $X$. All of the infinite paths which do not remain in $X$ end up in cycles and they are a countable set. The dimension of $E$ is equal to the dimension of the set which consists of points determined by the infinite words accepted by $X$ fixing a vertex as the initial state. Let $\mathrm{a}, \mathrm{b}$ and c be the states shown in Figure 5. Let $A=L(X, \mathrm{a}), B=L(X, \mathrm{~b}), C=L(X, \mathrm{c})$ be the sets of the infinite words accepted by $X$ with the initial satetes, $\mathrm{a}, \mathrm{b}, \mathrm{c}$, respectively. Then $A, B, C$ satisfy the following set-equations, namely we obtain a graph iterated function system [5].

$$
\left\{\begin{array}{l}
A=f_{1}(A) \cup f_{2}(A) \cup f_{3}(B)  \tag{1}\\
B=g_{1}(A) \cup g_{2}(B) \cup g_{3}(C) \\
C=h_{1}(A) \cup h_{2}(A) \cup h_{3}(C)
\end{array} .\right.
$$




Figure 4: decomposition to strongly connected components


Figure 5: $X$

$$
\begin{aligned}
& f_{1}(x)=\alpha^{4} x+\alpha^{2}+\alpha+1 \\
& f_{2}(x)=\alpha^{4} x+\alpha^{2}+\alpha \\
& f_{3}(x)=\alpha^{5} x+\alpha^{4}+\alpha^{2}+\alpha \\
& g_{1}(x)=\alpha^{7} x+\alpha^{5}+\alpha^{4}+\alpha^{3}+\alpha+1 \\
& g_{2}(x)=\alpha^{4} x+\alpha^{3}+\alpha \\
& g_{3}(x)=\alpha^{5} x+\alpha^{4}+\alpha^{3}+\alpha \\
& h_{1}(x)=\alpha^{10} x+\alpha^{8}+\alpha^{7}+\alpha^{6}+\alpha^{4}+\alpha^{3}+\alpha^{2}+1 \\
& h_{2}(x)=\alpha^{5} x+\alpha^{3}+\alpha^{2}+1 \\
& h_{3}(x)=\alpha^{4} x+\alpha^{3}+\alpha^{2}
\end{aligned}
$$

The Mauldin-Williams graph $G$ for this system is shown in Figure 6. The dimension $s$


Figure 6: Mauldin-Williams graph for $E$
associated to $G$ is computed as follows [5, 9]:

$$
\begin{gather*}
\operatorname{det}\left(\begin{array}{ccc}
2|\alpha|^{4 s}-1 & |\alpha|^{5 s} & 0 \\
|\alpha|^{7 s} & |\alpha|^{4 s}-1 & |\alpha|^{5 s} \\
|\alpha|^{5 s}+|\alpha|^{10 s} & 0 & |\alpha|^{4 s}-1
\end{array}\right)=0  \tag{2}\\
|\alpha|^{20 s}-|\alpha|^{16 s}+|\alpha|^{15 s}+3|\alpha|^{12 s}-5|\alpha|^{8 s}+4|\alpha|^{4 s}-1=0 .
\end{gather*}
$$

Thus $|\alpha|^{s}$ can take the following four real values,
-1 ,
$-0.9704028248572908964456 \cdots$,
$-0.8265233553562334820903 \cdots$,
$0.79254349837573255006271 \cdots$.

Since $s>0$, we obtain $|\alpha|^{s}=0.792543498375732550062 \cdots$ and $s=1.15931959819470279575$ $\cdots$. By using the result of [5], this value is the upper bound of the Hausdorff dimension of $A, B, C$.

Theorem 2 E has the Hausdorff dimension smaller than or equal to $1.15931959819470279575 \cdots$.

From numerical experiments, the upper bound in the theorem above seems to be exactly equal to the Hausdorff dimension of $E$.

Conjecture 1 Each of the sets, $f_{1}(A) \cap f_{2}(A), f_{2}(A) \cap f_{3}(B), f_{3}(B) \cap f_{1}(A), g_{1}(A) \cap g_{2}(B), g_{2}(B) \cap$ $g_{3}(C), g_{3}(C) \cap g_{1}(A), h_{1}(A) \cap h_{2}(A), h_{2}(A) \cap h_{3}(C), h_{3}(C) \cap h_{1}(A)$ consits of only one point.
$A, B$ and $C$ has the same Hausdorff dimension $s$, and each of the $s$-dimensional Hausdorff measure of $A, B, C$ is positive and finite. (See the first half of the proof of Corollary 3.5 in [6].) Regarding this conjecture to be true, we have

$$
\begin{aligned}
\mathcal{H}^{s}(A) & =\mathcal{H}^{s}\left(f_{1}(A)\right)+\mathcal{H}^{s}\left(f_{2}(A)\right)+\mathcal{H}^{s}\left(f_{3}(B)\right) \\
& =|\alpha|^{4 s} \mathcal{H}^{s}(A)+|\alpha|^{4 s} \mathcal{H}^{s}(A)+|\alpha|^{5 s} \mathcal{H}^{s}(B)
\end{aligned}
$$

where $\mathcal{H}^{s}(A)$ denotes the $s$-dimensional Hausdorff measure of $A$. In the same way, we obtain

$$
\left(\begin{array}{c}
\mathcal{H}^{s}(A) \\
\mathcal{H}^{s}(B) \\
\mathcal{H}^{s}(C)
\end{array}\right)=\left(\begin{array}{ccc}
2|\alpha|^{4 s} & |\alpha|^{5 s} & 0 \\
|\alpha|^{7 s} & |\alpha|^{4 s} & |\alpha|^{5 s} \\
|\alpha|^{5 s}+|\alpha|^{10 s} & 0 & |\alpha|^{4 s}
\end{array}\right)\left(\begin{array}{c}
\mathcal{H}^{s}(A) \\
\mathcal{H}^{s}(B) \\
\mathcal{H}^{s}(C)
\end{array}\right)
$$

and we have (2).
Figure 7 shows the points $\left\{z:(z, w) \in T(0) \cap T \cap T^{\prime}, T, T^{\prime} \in \mathcal{T}\right\}$, which seems to have the same dimension as that of $T(0) \cap T(1) \cap T(11)$.

3 Appendix The transision function of the automaton $M$ are shown below. The notation

$$
\begin{aligned}
& \mathrm{m} \\
& =0=>\mathrm{n} \\
& =0=>1 \\
& =1=>\mathrm{k}
\end{aligned}
$$

means that there are edges from the state $m$, one to the state $n$ labeled by 0 , one to $l$ labeled by 0 , and one to $k$ labeled by 1 .

$$
\begin{aligned}
& 1 \\
& =0=>2 \\
& =0=>106 \\
& =0=>109 \\
& =0=>111 \\
& =1=>114 \\
& 2 \\
& =0=>3 \\
& =0=>99 \\
& =0=>100 \\
& =1=>103 \\
& =1=>104 \\
& 3 \\
& =1=>4 \\
& 4 \\
& =1=>5 \\
& 5 \\
& =0=>6 \\
& 6 \\
& =0=>3 \\
& =0=>7 \\
& =0=>80 \\
& =1=>93 \\
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& =0=>16 \\
& 16 \\
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& =0=>18 \\
& =0=>20 \\
& =0=>61 \\
& =1=>64 \\
& =1=>77
\end{aligned}
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| 90 | 99 <br> $=0=>91$ | 108 <br> $=1=>14$ <br> 91 |
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| $=1=>92$ | 100 | 109 |
| 92 | $=0=>101$ | $=0=>110$ |
| $=0=>89$ | $=1=>102$ | 110 |
| 93 | 101 | $=1=>57$ |
| $=1=>94$ | $=1=>82$ | 107 |
| 94 | 102 | $=1=>108$ |
| $=1=>95$ | $=1=>88$ | 112 |
| 95 | 103 | $=1=>113$ |
| $=0=>37$ | $=1=>94$ | 113 |
| 96 | 104 | $=1=>30$ |
| $=1=>97$ | $=1=>105$ | 114 |
| 97 | 105 | $=0=>3$ |
| $=1=>98$ | $=1=>98$ | 111 |
| 98 | 106 | $=0=>112$ |
| $=0=>6$ | $=0=>107$ |  |

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Department of Administration,
Prefectural University of Kumamoto, Tsukide 3-1-100, Kumamoto,
Japan 862-8502
E-mail; sadahiro@pu-kumamoto.ac.jp

